# Skipless Chain Decompositions & Improved Poset Saturation Bounds

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- relation:  $\subseteq$



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True for every poset, and every way to embed it.

#### Structural Theorem [B., Groenland, Jacob, Johnston, 2022+]

Any subposet  $\mathcal{P}$  of  $2^{[n]}$  with largest antichain of size k can be **covered** by a family of k **disjoint skipless** chains in  $2^{[n]}$ .

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We generalise a result of Lehman and Ron (2001) who proved the special case where all chains of the family are of size 2 and all top (resp. bottom) elements of the chain have the same size.

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 $<sup>{}^{1}\</sup>mathcal{F} \subseteq 2^{[n]}$  is convex if for all  $X, Z \in \mathcal{F}$  and  $X \subset Y \subset Z, Y \in \mathcal{F}$ .

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## Sketch of the sketch of the proof

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Any family of k chains in  $2^{[n]}$  can be **covered** by a family of k **disjoint skipless** chains in  $2^{[n]}$ .



# $\mathcal{F} \subseteq 2^{[n]}$ , is *k*-saturated if:

- $\mathcal{F}$  has no antichain of size k;
- $\mathcal{F} \cup \{x\}$  has an antichain of size k for any  $x \in 2^{[n]} \setminus \mathcal{F}$ .

 $sat^*(n, k) = minimum |\mathcal{F}|$  over all k-saturated families  $\mathcal{F}$  in  $2^{[n]}$ .

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Red sets form an 2-saturated family for the hypercube  $2^{[3]}$ : sat\*(3,2)  $\leq$  4. Can we extend this construction to *k*-saturated ?



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Conjecture (FKKMRSS):  $\forall k \geq 2$ , sat\* $(n, k) \sim n(k-1)$  as  $n \to \infty$ .



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There is a half page proof of both conjecture with our decomposition lemma!

 $\nu(\mathcal{F}) \rightarrow$  the size of the maximum matching from  $\mathcal{F}$  to its shadow  $\partial \mathcal{F}$ .  $\mathcal{C}(m, t) \rightarrow$  initial segment of colex of size *m* on layer *t*.

Define the sequence  $c_{\lfloor \ell/2 \rfloor} = k - 1$ , and for  $0 \le t < \lfloor \ell/2 \rfloor$ , let  $c_t = \nu \left( \mathcal{C}(c_{t+1}, t+1) \right)$ .

B, Groenland, Jacob and Johnston (2023+) For  $n \ge 2 \log(k) + 1$ ,

$$\operatorname{sat}^*(n,k) = 2\sum_{t=0}^{\lfloor \ell/2 
floor} c_t + (k-1)(n-1-2\lfloor \ell/2 
floor).$$

The lower bound still holds for  $n \ge \ell$  (and sat<sup>\*</sup> $(n, k) = 2^n$  for  $n < \ell$ ).

#### Definition

- $\mathcal{F} \subseteq 2^{[n]}$  a set system is  $\mathcal{P}$ -saturated if:
  - $\mathcal{F}$  has no induced copy of  $\mathcal{P}$ ;
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Theorem (Morrison, Noel and Scott 2014;

 $\leq$  sat\* $(n, C_k) \leq 2^{0.98k}$ 

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Theorem (Morrison, Noel and Scott 2014; Gerbner, Keszegh, Lemons, Palmer, Pálvölgyi, Patkós 2011)

 $2^{(k-3)/2} \leq \mathsf{sat}^*(n, C_k) \leq 2^{0.98k}$ 

# Table

$\mathbf{poset} \ P$	$\mathbf{sat}(n, P)$	$\mathbf{sat}^*(n, P)$	
$C_2$ , chain	= 1	= 1	
$A_2$ , antichain	= 1	= n + 1	
$C_3$ , chain	= 2	=2	
$C_2 + C_1$ , chain and single	= 2	= 4	case analysis
$\vee$ fork (or $\wedge$ )	= 2	= n + 1	[F7]
$A_3$ , antichain	= 2	= 3n - 1	[F7]
$C_4$ , chain	= 4	= 4	[G6]
$\vee_3$ , fork with three times	= 3	$\geq \log_2 n$	[F7]
$\diamond$ , diamond	= 3	$\geq \sqrt{n}$	[MSW]
		$\leq n+1$	[F7]
$\diamond^-$ , diamond minus an edge	= 3	= 4	case analysis
$\bowtie$ , butterfly	= 4	$\geq n+1$	[I]
		$\leq 6n - 10$	[Thm  3.16]
Y	= 3	$\geq \log_2 n$	[Thm. <u>3.6</u> ]
N	= 3	$\geq \sqrt{n}$	[I]
		$\leq 2n$	[F7]
$2C_2$	= 3	$\geq n+2$	[Thm. 3.11]
		$\leq 2n$	[Prop. <u>3.9</u> ]

#### Figure 1: Table from Keszegh, Lemons, Martin, Pálvölgyi and Patkós 2022

### Table

$C_3 + C_1$ , chain and single	= 3	$\leq 8$	[Prop. 3.18]
$\vee +1$ , fork and single	= 3	$\geq \log_2 n$	[F7]
$C_2 + A_2$	= 3	$\leq 8$	[Prop. 3.18]
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		$\leq 4n+2$	[F7]
$C_5$ , chain	= 8	= 8	[G6]+[MNS]
$C_6$ , chain	= 16	= 16	[G6]+[MNS]
$C_k$ , chain $(k \ge 7)$	$\geq 2^{(k-3)/2}$	$\geq 2^{(k-3)/2}$	[G6]
	$\leq 2^{0.98k}$	$\leq 2^{0.98k}$	[MNS]
$A_k$ , antichain	= k - 1	$\geq \left(1 - rac{1}{\log_2 k} ight) rac{k}{\log_2 k} n$	[MSW]
		$\leq kn - k - \frac{1}{2}\log_2 k + O(1)$	[F7]
$3C_2$	= 5	$\leq 14$	[Prop. <u>3.13</u> ]
$5C_2$	= 9	$\leq 42$	[Prop. 3.18]
$7C_2$	= 13	$\leq 60$	[Prop. 3.18]
any poset on $k$ elements	$\leq 2^{k-2}$		[Thm. 1.1]
UCTP (def. in Section $3.2$ )	O(1)	$\geq \log_2 n$	[F7]
UCTP with top chain	O(1)	$\geq \log_2 n$	[Thm. 3.6]
chain + shallower	O(1)	O(1)	[Thm. 3.8]

Figure 2: Table from Keszegh, Lemons, Martin, Pálvölgyi and Patkós 2022

Very recently, a general lower bound has been shown.

Theorem (Freschi, Piga, Sharifzadeh and Treglown 2023)

For any poset P either sat\* $(n, P) \ge 2\sqrt{n} - 2$  or sat\* $(n, P) = O_P(1)$ .

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Theorem (B., Groenland, Ivan, Johnston, 2023+)

For any poset P, sat\* $(n, P) \leq n^{|P|^2}$ .

For every poset  $\mathcal{P}$ , either sat\* $(n, \mathcal{P}) = O_{\mathcal{P}}(1)$  or sat\* $(n, \mathcal{P}) = \Theta_{\mathcal{P}}(n)$ .

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Thank you!

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